A Data-Driven Frequency Scaling Approach for Deadline-aware Energy Efficient Scheduling on Graphics Processing Units (GPUs)

Shashikant Ilager, Rajeev Muralidhar, Kotagiri Rammohanrao, Rajkumar Buyya  
Cloud Computing and Distributed Systems (CLOUDS) Laboratory  
School of Computing and Information Systems, The University of Melbourne, Australia

Abstract—Modern computing paradigms, such as cloud computing, are increasingly adopting GPUs to boost their computing capabilities primarily due to the heterogeneous nature of AI/ML/deep learning workloads. However, the energy consumption of GPUs is a critical problem. Dynamic Voltage Frequency Scaling (DVFS) is a widely used technique to reduce the dynamic power of GPUs. Yet, configuring the optimal clock frequency for essential performance requirements is a non-trivial task due to the complex nonlinear relationship between the application’s runtime performance characteristics, energy, and execution time. It becomes more challenging when different applications behave distinctively with similar clock settings. Simple analytical solutions and standard GPU frequency scaling heuristics fail to capture these intricacies and scale the frequencies appropriately. In this regard, we propose a data-driven frequency scaling technique by predicting the power and execution time of a given application over different clock settings. We collect the data from application profiling and train the models to predict the outcome accurately. The proposed solution is generic and can be easily extended to different kinds of workloads and GPU architectures. Furthermore, using this frequency scaling by prediction models, we present a deadline-aware application scheduling algorithm to reduce energy consumption while simultaneously meeting their deadlines. We conduct real extensive experiments on NVIDIA GPUs using several benchmark applications. The experiment results have shown that our prediction models have high accuracy with the average RMSE values of 0.38 and 0.05 for energy and time prediction, respectively. Also, the scheduling algorithm consumes 15.07% less energy as compared to the baseline policies.

Index Terms—GPU, Energy, Data-Driven, Scheduling, Machine Learning.

I. INTRODUCTION

Graphics Processing Units (GPUs) have become ubiquitous in modern computing paradigms and platforms, such as Cloud computing and supercomputing environments, due to their massive computational capabilities. Furthermore, the Single Instruction Multiple Data (SIMD) architecture of GPUs is ideally suitable for many parallel and compute-intensive scientific and business workloads [1], [2]. These advantages manifested into the deployment of a large number of GPU clusters in many data centers, including Top500 supercomputers and also in the public Clouds [3], [4], [5], [6]. In spite of this increased usage, the power consumption of GPUs has become a significant bottleneck for designing hyper-scale GPU systems [7], [8]. On the other hand, GPU workloads are more sensitive to their Quality of Service (QoS) constraints requiring faster execution and thus spending more energy. Therefore, energy-efficient workload management with QoS satisfaction is exceedingly essential.

Dynamic Voltage Frequency Scaling (DVFS) is a popular technique to reduce active power by varying the GPU frequencies [9], [10], [11]. The modern GPUs have two frequency domains, core and memory, each with many numbers of frequency ranges. While former regulates the Streaming Multiprocessors (SM) (including register, texture cache, shared memory, and L2 cache), and the latter governs bandwidth of DRAM [11]. For instance, NVIDIA Tesla P100 GPU supports one memory frequency (715 MHz) and 62 core frequencies ([544-1328] MHz), and NVIDIA GTX 980 supports four memory frequencies ([3505-324] MHz) and 87 different core frequencies([135-1428] MHz ) with the total number of 267 possible frequency combinations. A particular combination of memory and core frequency can be set using the NVIDIA Management Library (NVML). However, the principle DVFS notion—higher frequency range increases the performance requiring more power, while lower frequency consumes less power by decreasing the performance do not hold in all the scenarios [12]. In addition, different GPU application kernels behave differently concerning energy and performance with the frequency settings due to their different resource footprints and the intricate instruction execution patterns. Thus, due to such non-linear dependencies, estimating and optimally scaling the frequencies for a given application is non-trivial.

Furthermore, frequency scaling becomes more challenging when a scheduler needs to schedule multiple applications with their deadline requirements. In such a case, the scheduler should not only identify the energy-efficient frequency combinations, but it also needs to take care of the application’s execution time. Such scenarios are highly prevalent in real-time HPC and cloud environments [13], [14]. However, existing analytical and heuristic-based GPU frequency scaling [15], [16] methods are inefficient as they fail to capture the complex non-linearity between the frequency settings, performance, and power. To that end, data-driven DVFS scaling is a promising technique that is built using actual measurements. Models built using such methods can accurately scale the frequencies based on application demands[17]. Moreover, once the model is trained, the new applications can be scheduled on-the-fly with minimum profiling data.
In this paper, we present a data-driven approach for frequency scaling by observing key architectural, power, and performance counters and predicting the estimated application power and execution time. In addition, guided by these prediction models, we propose a deadline-aware energy-efficient scheduling algorithm that accurately scales the GPU frequency according to the application requirements. We use twelve applications for evaluation from two standard GPU benchmarking suites, Rodinia [18] and Polybench [19]. The training data is generated from profiling the applications using nvprof, a standard profiling tool from NVIDIA. Furthermore, several machine learning models are explored to accurately predict the energy and execution time of applications for the given frequency domains. Based on the experimental results, CatBoost, an ensemble-based gradient boosting learning model, is chosen for prediction modeling. We implement the prototype scheduling system and evaluate the proposed techniques on real platforms. The experimental results conducted on the NVIDIA GPU device, Tesla P100 (Pascal micro-architecture), have shown that our prediction models have high accuracy and proposed scheduling algorithm consumes less energy as compared to the baseline algorithms.

In summary, we make the following key contributions:

- We propose a data-driven prediction model to accurately predict the energy and execution time of applications to assist the efficient frequency scaling configuration by observing key architectural, power, and performance counters and metrics.
- We design and present a deadline-aware energy-efficient application scheduling algorithm using the prediction models.
- We implement a prototype system and evaluate the proposed solution on a real platform using standard benchmarking applications.
- We show the efficiency and efficacy of our proposed solution with extensive experiments, and results are compared and analyzed with the existing state-of-the-art solutions.

The remainder of the paper is organized as follows. Section II describes the DVFS background, motivation of the work, and system model. Section III presents the data-driven frequency scaling techniques. Section IV shows our proposed deadline aware energy-efficient scheduling algorithm. Section V describes the performance evaluation with the analysis of the results. The related work is explained in Section VI. Finally, Section VII draws the conclusions and future work.

II. BACKGROUND MOTIVATION AND SYSTEM MODEL

A. GPU DVFS

The power consumption of a GPU (or any semiconductor logic block in general) is a combination of dynamic \( P_{\text{dynamic}} \) and static power \( P_{\text{static}} \). The static power \( P_{\text{static}} \) is related to the leakage and energy consumed when the system is idle, and usually it is managed by hardware and/or software using different sleep states [20]. However, large amounts of energy is spent on the dynamic power \( P_{\text{dynamic}} \) which is proportional to the run time of the workload. Performance management of GPU typically rely on the DVFS-based heuristics to regulate the dynamic power to save the energy. The frequency is normally regulated based on the application’s utilisation parameters or system’s temperature threshold throttling mechanisms [17]. Thus, the dynamic power can defined as below:

\[
P_{\text{dynamic}} \propto V^2 F
\]

In Equation 1, \( F \) denotes the operating frequency, while \( V \) denotes the supply voltage. Based on the current operating frequency, a combination of hardware and software changes the frequency (and thereby the underlying voltage); certain frequency ranges can share the same voltage level. Furthermore, GPUs have multiple frequency domains \( F = \{ f_{\text{memclock}}, f_{\text{memclock}} \} \), regulating hardware components related to streaming multiprocessor or graphics processor and the memory components, respectively [12]. Thus, considering that, usually, hardware logic manages the voltage based on operating frequency, we focus on benefiting from regulating frequency and scaling it based on application behaviors.

B. Motivation

Estimating the optimal frequency is a non-trivial problem due to the complex behaviors of applications regarding their energy consumption and execution time. To analyze this complexity, we plot the behavior of different applications towards energy and execution time by changing the core frequency of the GPU, as shown in Figure 1. These executions are from NVIDIA Tesla P100 GPU that has one memory frequency and 62 core frequencies. So we only vary the core frequencies. In Figure 1, we can observe that when the core frequency increases, energy consumption is not always linear. And also,

![Fig. 1: Power, time and clock relationship of different applications](image-url)
the lavaMD (Figure 1a) application has a completely inconsistent response to frequency variations. Furthermore, some application produces different functionalities between certain frequency range, for application CORR (Figure 1b), energy consumption has a non-convex curve between [730-920] MHz. Similar nonlinear behavior is present in Figure 1c and 1d, where execution time and energy consumption have some unexpected spikes and dips in their energy response. Moreover, optimising such non-linear non-convex functions is an NP-hard problem [21]. Therefore, it is extremely challenging to find energy-efficient frequency combinations under application’s QoS constraints. Simple analytical models that linearly regulate the core frequencies are inaccurate and inefficient to reduce the energy or increase the performance [17]. Motivated by these factors, we model the frequency scaling problem based on the data-driven methods.

C. System Model

A high-level system model of the proposed framework is shown in Figure 2. It is broadly classified into two parts, predictive modeling and data-driven scheduler. In the first part, we collect the training data that consists of three parts, profiling information, energy-time measurements, and respective frequency configurations. We then predict two entities for a given application and frequency configuration, i.e., energy consumption and execution time. Afterward, in the second part, the new applications arrive with the deadline requirements and minimize profiling data from a default clock frequency execution. The scheduler finds correlated application data using the clustering technique, and this data is used for predicting the energy and execution time over all frequencies. Finally, based on the deadline requirements and energy efficiency, the scheduler scales the frequencies and executes the applications. The important components of this framework are explained in the following sections.

III. DATA-DRIVEN FREQUENCY SCAILING FOR GPUs

In this section, we discuss the prediction models in detail, which include data collection, training, and model evaluation.

A. Feature Collections

ML-based models are trained using real-time measurement data from the environment, and these models are used in the run time to predict the outcomes. In case of GPU energy and performance prediction, several existing studies rely either on static source code metrics [16], [22], or on run time traces and profiling data [23]. The profiling based method is most suitable due to its ability to gather the real resource footprints and hardware counters of applications on particular GPU architectures, which is crucial to estimate energy and time accurately.

The input to our training model consists of (1) applications profiled features \( F = \{f_1, f_2, f_3,... f_n\} \), (2) respective GPU frequency pair \( \{\text{smclock}, \text{memclock}\} \), and (3) energy and execution time measurements. The profiling features contain the information of an application’s run time metrics related to its different hardware components utilization values, instruction counts, communication, and cache metrics, etc.

Applications: We have considered twelve different applications from two heterogeneous computing benchmark suites, Polybench [19] and Rodinia [18]. These two benchmark suites cover a wide range of applications. The Polybench suit covers many linear algebraic applications while the Rodinia suit covers different scientific applications. The details of these applications, including their domain and input specifications, are shown in Table I.

Profiling: For profiling the applications, we use nvprof, a standard profiling tool by NVIDIA for CUDA applications. Although NVIDIA has recently released new nsight-systems toolkits for monitoring and profiling, they do not support many existing GPU architectures and CUDA versions, so we use nvprof. We have also used the nvidia-smi toolkit, which is built on top of nvml library, a C-based API for monitoring and managing various states of the NVIDIA GPU devices. This tool allows application users to set the supported GPU application clocks and also to measure the energy consumption metrics.

We have developed Python scripts to collect the profiling metrics that runs all the applications iteratively on different frequency domains supported by GPU. Initially, we collect all the available metrics provided by nvprof using --metrics all argument and export the output in csv format. The energy and execution time are gathered by running applications separately to avoid the effect of profiling on these metrics.

We collect metrics from every alternative clock pair of the Tesla P100 GPU from its supporting 62 combinations of core and memory frequencies to reduce the data collection time. It is important to note that, some applications take up to ten minutes for each profiling session as nvprof replays the application kernels over several passes to collect the metrics. The nvprof provides more than 120 metrics of GPU counters for each execution. For the sake of brevity, we list the top twenty features that dominate in energy and execution time prediction in Table II. The details about selecting these features are explained in feature analysis Section III-C. Here, the features sm (SM’s utilisation level) is collected from nvidia-smi dmon API and remaining all are from nvprof.

These collected application profiling metrics, along with the frequency configuration, energy, and time metrics, constitute
the total training data, which are then used to build predictive models.

### B. Prediction Models

When building any predictive models, it is often required to test the suitable candidate algorithms and adopt the model that works best for the given input training data and the problem domain. In this regard, the prediction of energy and execution time requires regression-based machine learning (ML) models. We have investigated several suitable candidate ML algorithms, including Linear Regression (LR), lasso-linear regression (Lasso), and Support Vector Regression (SVR). Also, we explored ensemble-based gradient boosting techniques, eXtreme Gradient Boosting (XGBoost), and CatBoost. The goal is to build energy and execution time prediction models for each GPU device to assist the frequency configuration. To that end, prediction models are trained for two outputs, i.e., energy ($E$) and execution time ($T$). The profiling data from all the applications are partitioned into training and testing datasets with 70% and 30%, respectively.

The input feature set now includes a set of tuples with each tuple having profiled features and frequency combination i.e., $F = \{f_1, f_2, \ldots, f_n\} \cup \{f_{\text{smclock}}, f_{\text{memclock}}\}$ while the two models output predicted energy consumption $P$ and execution time $T$.

We use the sci-kit learn package [22] to implement the LR, Lasso, and SVR. For XGBoost and CatBoost, the standard python packages are used that are publicly available \(^1\) \(^2\). A few of the profiling parameters from `nvprof` are categorical, representing different components utilization as low, mid, and high. Among a total of 120 features collected, 15 features were categorical including `dram_utilisation`, `double_precision_fu_utilisation`. Only numerical features are used in the models except for CatBoost. However, CatBoost is specifically designed to work with the categorical or mixed data, and it has an efficient way of representing the categorical variables. Here, the categorical features are transformed into numerical features by the technique of order target statistics.

The parameters for each of the algorithms are the default parameters, and self-explanatory in our implementation. To avoid overfitting of the models, we adopt the leave-one-out cross-validation, where we exclude the data from a particular application in training and evaluate this model with the excluded application’s test data. This helps to assess the robustness of models, and proven efficiency will help to use these offline trained models for new applications without retraining the models.

The goodness of fit is measured using the Root Mean Square Error (RMSE) metric, which is a standard evaluation metric in regression-based problems [23]. The RMSE is defined as:

$$RMSE = \sqrt{\frac{1}{n} \sum_{i=0}^{n} (y_i - \hat{y}_i)^2}$$

In Equation 2, $y_i$ is the observed value, $\hat{y}_i$ is the predicted output variable by prediction model, and $n$ is the total number of predictions. The value of RMSE represents the standard deviation of the residuals or prediction errors. It also indicates how far are the data points from the model fitted line. Thus, lower RMSE values are preferred.

\(^1\)https://catboost.ai/
\(^2\)https://github.com/dmlc/xgboost

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### TABLE I: Description of applications

<table>
<thead>
<tr>
<th>Application</th>
<th>Domain/Description</th>
<th>Suite</th>
<th>Input</th>
</tr>
</thead>
<tbody>
<tr>
<td>particlefilter_naive</td>
<td>Medical Imaging</td>
<td>Rodinia</td>
<td>-x 128 -y 128 -z 10 -np 1000</td>
</tr>
<tr>
<td>particlefilter_float</td>
<td>Medical Imaging</td>
<td>Rodinia</td>
<td>-x 128 -y 128 -z 10 -np 1000</td>
</tr>
<tr>
<td>myocyte</td>
<td>Biological Simulation</td>
<td>Rodinia</td>
<td>10000, 1000, 1</td>
</tr>
<tr>
<td>lavaMD</td>
<td>Molecular Dynamics</td>
<td>Rodinia</td>
<td>-boxes 1d 50</td>
</tr>
<tr>
<td>Backprop</td>
<td>Pattern Recognition</td>
<td>Rodinia</td>
<td>983040</td>
</tr>
<tr>
<td>SYR</td>
<td>Symmetric rankk operations</td>
<td>Polybench</td>
<td>M 1024, N 1024</td>
</tr>
<tr>
<td>SYR2K</td>
<td>Symmetric rankk2 operations</td>
<td>Polybench</td>
<td>M 2048, N 2048</td>
</tr>
<tr>
<td>GEMM</td>
<td>Matrix Multiply C = A x B + C</td>
<td>Polybench</td>
<td>N 2048, NJ 2048, NK 2048</td>
</tr>
<tr>
<td>COVAR</td>
<td>Covariance Computation</td>
<td>Polybench</td>
<td>M 2048, N 2048</td>
</tr>
<tr>
<td>CORR</td>
<td>Correlation Computation</td>
<td>Polybench</td>
<td>M 2048, N 2048</td>
</tr>
<tr>
<td>ATAX</td>
<td>Matrix Transpose and Vector Multiplication</td>
<td>Polybench</td>
<td>NX 16384, NY 16384</td>
</tr>
<tr>
<td>2MM</td>
<td>2 Matrix Multiplications (D=A.B; E=C.D)</td>
<td>Polybench</td>
<td>NI 4096, NJ 4096, NK 4096, NL 4096</td>
</tr>
</tbody>
</table>

### TABLE II: Features used in energy and time prediction (top 20)

<table>
<thead>
<tr>
<th>Power</th>
<th>Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>sm</td>
<td>sm</td>
</tr>
<tr>
<td>sm_clock</td>
<td>l2_t tex_read_hit_rate</td>
</tr>
<tr>
<td>l2_t tex_read_hit_rate</td>
<td>l2_t tex_read_transactions</td>
</tr>
<tr>
<td>tex_cache_throughput</td>
<td>tex_cache_transactions</td>
</tr>
<tr>
<td>ipc</td>
<td>dram_write_transactions</td>
</tr>
<tr>
<td>flop_dp_efficiency</td>
<td>ipc</td>
</tr>
<tr>
<td>shared_load_throughput</td>
<td>inst_executed_shared_loads</td>
</tr>
<tr>
<td>stall_exec_dependency</td>
<td>gt_efficiency</td>
</tr>
<tr>
<td>stall_inst_fetch</td>
<td>inst_reply_overhead</td>
</tr>
<tr>
<td>eligible_warp_per_cycle</td>
<td>inst_executed_shared_stores</td>
</tr>
<tr>
<td>stall_constant_memory_dependency</td>
<td>l2_read_throughput</td>
</tr>
<tr>
<td>pcie_total_data_transmitted</td>
<td>dram_throughput</td>
</tr>
<tr>
<td>dram_read_transactions</td>
<td>dram_execution_efficiency</td>
</tr>
<tr>
<td>dram_read_bytes</td>
<td>dram_read_bytes</td>
</tr>
<tr>
<td>issue_slots</td>
<td>local_store_throughput</td>
</tr>
<tr>
<td>l2_t tex_write_throughput</td>
<td>glb_efficiency</td>
</tr>
<tr>
<td>inst_bit_convert</td>
<td>global_store_requests</td>
</tr>
<tr>
<td>l2_global_load_bytes</td>
<td>stall_memory_throttle</td>
</tr>
<tr>
<td>glb_requested_throughput</td>
<td>dram_utilisation</td>
</tr>
<tr>
<td>pcie_total_data_received</td>
<td>inst_fp_32</td>
</tr>
</tbody>
</table>
The performance of different algorithms is shown in Figure 3. Here, we can observe that the CatBoost has the lowest RMSE value of 0.38 in energy prediction, indicating residuals or prediction errors are less, and its predictions are more accurate. We observed that linear models like LR, SVR, and Lasso perform worst in estimating energy and slightly better in predicting the execution time. It is because energy consumption has more non-linearity with the input features than the execution time, and simple linear models do not perform well at it. While in execution time prediction (Figure 3b), both XGBoost and CatBoost has equal performance (RMSE value of 0.05). As the performance of the CatBoost is promising in both models, we choose it as our prediction algorithm.

We perform hyperparameter tuning to further optimize the CatBoost model; we use the grid search technique over the parameter space. The results of the grid search are shown in Table III. Here, the parameters iterations and depth decides the number and size of the decision trees while learning_rate is used for reducing the gradient step. The $l_2\_leaf\_reg$ represents the coefficient at the L2 regularisation term of the cost function.

### C. Feature Analysis

We carry feature analysis to understand the importance of individual features towards the performance of the prediction model. It also represents the features that are highly influential on the prediction output.

Figure 4 indicates the Feature Importance ($F.I$) score of different features. We plot the twenty most significant features sorted in descending order of their score. Here, $F.I$ value represents the difference between the loss value of the model with and without that feature. The model without this feature is similar to the one that would have been trained if this feature was excluded from the data set. Since, RMSE is our loss function, the F.I score on y axis shows change in RMSE value.

We can observe from Figure 4a and Figure 4b, different types of features contribute to different magnitude while predicting energy and time of application, respectively. In both cases, feature smi, which represents the streaming multi-processor’s utilization, has the highest $F.I$ score showing its high impact on the energy and time. Furthermore, sm_clock is the second most important feature in predicting energy, reflecting a direct correlation between energy and frequency clock settings. Please note, since our testbed GPUs (Tesla-P100) have only one memory frequency ($f_{memclock}$), it does not feature in the top twenty features as there is no variation introduced by it in the data set. For the system with multiple $f_{memclock}$, it is expected that it would have a significant effect on the model’s performance. We can also observe from time model in Figure 4b, different features present when compared to the energy model (Figure 4a). A total of 5 features are in common between two models. The features related to l2 cache and stall dependencies have more impact in the energy model, while in time prediction, separate features like inst_executed_shared_stores, inst_fp_32 have occurred in top 20 features showing the higher co-relation between execution time and the metrics related to instruction count.

To further analyze the impact of the number of features on the prediction model’s performance, we carry a threshold analysis. Here, features are sorted based on their $F.I$ score, and recursively added to the training data set, and resulting loss value (RMSE) is calculated accordingly. As shown in the Figure 5, for both the power and time model, the top 20 features are sufficient enough to achieve the reasonable performance with excellent RMSE value and further inclusion of features do not yield much improvement in the result without increasing the training cost.

### D. Feature Correlation with Clustering

The prediction models need exhaustive application profiling data from multiple frequency combinations. However, profiling every new application is tedious and infeasible in real-time. Thus, using the existing data and correlating with the new application is a common practice in profiling-based predictive modeling scenarios [24]. In such a case, a new application should have at least one set of profiling data of one frequency combination; we take the default clock as a reference for this. We generate the clusters based on already collected exhaustive data and predict the cluster label for a new application. Furthermore, a highly correlated application within the cluster is chosen from simple heuristic, i.e., application with the

<table>
<thead>
<tr>
<th>Catboost model</th>
<th>depth</th>
<th>l2_leaf_reg</th>
<th>iterations</th>
<th>learning_rate</th>
</tr>
</thead>
<tbody>
<tr>
<td>Power</td>
<td>4</td>
<td>5</td>
<td>1200</td>
<td>0.1</td>
</tr>
<tr>
<td>Time</td>
<td>4</td>
<td>3</td>
<td>1200</td>
<td>0.03</td>
</tr>
</tbody>
</table>

TABLE III: Optimal parameters obtained for CatBoost from grid search technique.
lowest absolute difference in execution time is selected to further match the applications performance-similarity.

We use the same set of twelve applications to perform this analysis. To generate the clusters, we use K-means. We found that an optimal number of clusters is five based on the weighted sum of the squared error metric. Applications belonging to a different group and their correlated application can be seen in Table IV. The cluster-3 has only one application, i.e., 2MM, suggesting the essence of having a still more number of applications in the sample space to have at least two or more applications in each cluster. The robustness of this method is evaluated by predicting the execution time and energy for all the applications using the profile data of corresponding correlated applications. The average RMSE value of 3.19 and 1.11 is achieved for energy and time prediction, respectively, proving the feasibility of this method.

### IV. Deadline Aware Application Scheduling by Data-Driven DVFS

The advantages of power and performance estimations of GPU workloads are manifold. It is used in resource management techniques like scheduling [14], power capping [25], and also in the analysis of performance bottlenecks of workloads [10]. In this work, we propose deadline aware energy-efficient application scheduling guided by the data-driven DVFS.

The workload model of our scheduling is shown in Figure 6a. It consists of a set of applications represented as a vector $W = \{A_1, A_2, A_3...A_n\}$, with their own arrival and deadline times, represented as vectors $A = \{a_1, a_2, a_3,..., a_n\}$ and $D = \{d_1, d_2, d_3,..., d_n\}$ where $\forall a_i \in A$ and $\forall d_i \in D$, respectively.

As illustrated in Figure 6, the power curve for individual applications are non-linear with their execution time. The objective is to configure the frequency that meets application $A_i$’s deadline $d_i$ and also has the lowest energy consumption according to it’s power curve. Therefore, considering the energy consumption of application $i$ is $P_i$, then reducing the overall GPU energy consumption is formulated as a minimisation problem as below:

$$\text{minimize} \quad P_{total} = \sum_i^n P_i$$

subject to $\quad \forall T_i \leq d_i$ (3)

In the above Equation 3, $T_i$ is application’s execution time. The objective function minimises the total GPU energy consumption and the constraint makes sure that application is executed within the deadline. However, solving Equation 3 is equivalent to constrained global optimization which is an NP-hard problem [26]. As it is impractical to find an optimal solution in real-time, we present a heuristic algorithm in order to reduce the problem complexity suitable for on-line
Algorithm 1 Deadline-aware Scheduling by Data-Driven DVFS

Inputs:
1: $W$: list of applications to be executed (workload)
2: $D$: list of application’s deadline (d)
Output: GPU application clock and schedule
while true do
1. jobQue ← getCurrentArrivalJobs($W$)
2. jobQue ← sort(jobQue, key=D, order=asc)
3. clockList ← getGPUSupportedClocks(deviceID)
4. for each job in jobQue do
5. predictionInput ← getCorrelatedData(job)
6. minPower ← MAX
7. maxTime ← job.d
8. ($f_{smclock}$, $f_{memclock}$) ← NULL
9. for each clockSet in clockList do
10. $P_{job}$ ← PREDICTPWR(predictionInput, clockSet)
11. $T_{job}$ ← PREDICTTIME(predictionInput, clockSet)
12. if $P_{job}$ < minPower and $T_{job}$ < maxTime then
13. minPower ← $P_{job}$
14. maxTime ← $T_{job}$
15. ($f_{smclock}$, $f_{memclock}$) ← clockSet
16. end if
17. end for
18. SETGPUAPPLICATIONCLOCK($f_{smclock}$, $f_{memclock}$)
19. if not ($f_{smclock}$, $f_{memclock}$) NULL then
20. EXECUTE(job)
21. end if
22. end for
23. end while
End

The Algorithm 1 shows proposed scheduling algorithm for deadline-aware application scheduling with the data-driven DVFS. Its objective is to reduce the energy while meeting the application’s deadline and it is achieved by generating an efficient schedule sequence and also suitably scaling the GPU frequencies. The algorithm takes the application list, and their corresponding deadlines as an input and outputs suitable predicted clock and scheduling sequence for applications.

First, according to applications arrival time, the available jobs are sorted based on the deadlines in ascending order (line 4,5) to make sure the jobs with the earliest deadline are executed first. Considering the new arrived application has only default clock input profile data, we find its correlated application and use it’s exhaustive profile data for prediction (as explained in Section III-D). Furthermore, the power and execution time is predicted for all the supported GPU frequency clock sets (lines 12-14) using the prediction models proposed in Section IV. For a given job, the clocks which have the lowest power consumption and also the predicted execution time less than its deadline is selected (15-18). Finally, the selected application clock is configured, and the application is executed.

The time complexity of Algorithm 1 is polynomial. Assuming we have $n$ jobs to be scheduled on a GPU with $c$ number of clocks. The sorting of jobs require worst-case time complexity of $n \log n$ (line 5). Furthermore, each job has to be evaluated on all clock-sets and executed, that has a time complexity of $nc$ (line 7-20). Hence, the total complexity will be $(n \log n + nc)$. Considering $c$ is a constant for any given GPU, the Algorithm 1 has polynomial time complexity of $O(n\log n)$.

V. PERFORMANCE EVALUATION

We discuss the implementation of our proposed algorithm integrated with the prediction models. We also analyze and discuss the results compared to baseline algorithms.

A. Implementation

We implemented the proposed scheduling framework using Python language. We developed a multithreaded application where the main thread executes the algorithm 1’s logic and invokes the application execution files. Additionally, it also launches two other background threads, one to collect the GPU data related to energy metrics (by running bash scripts with nvidia – smi dmon) and other kills the background thread once the application execution is done. Furthermore, the application clocks are predicted based on the proposed model in Section IV, and these predicted clocks are set before executing the scheduled application using the NVML’s nvidia – smi’s API.

B. Experimental Setup

We use Grid’5000 testbed [27] for our experiments. It is a large-scale flexible testbed for experiment-driven research, specifically, designed for experimental evaluations of the energy-efficient techniques [27]. We have used Tesla-P100 GPUs for our experiments. This machine has a dual CPU Intel Xeon Gold 6126 processor with 12 cores per CPU and 192 GB of primary memory. The GPU itself has 3584 cores with 16 GB primary memory. The machines are deployed with Debian 10 as the operating system installed with CUDA 10.1 drivers.

C. Benchmarking Applications

We have used twelve applications from two bench-marking suits (PolyBench and Rodinia), which are also part of our predictive modeling (Table I). To formally produce the workload
model described in Figure 6a, initially, we use GPU default application clocks run time ([715, 1189] MHz for Tesla-P100) as a reference to our application’s execution time. Also, the application’s arrival time and deadline are calculated based on the normal distribution. For the arrival time, the minimum and maximum value range of distribution are set to (1, 50), and for the deadline, it is set to (1, 2× execution time). This is to make sure that the application’s deadline can have a maximum value of twice their execution time. All these applications are CUDA-based implementation, and configurations are shown in Table I.

D. Analysis of Results

We evaluate the proposed algorithm 1 against two baselines. (1): Default Clock (DC): GPU frequencies are set to default application clocks. The applications usually run on default clocks under normal conditions (2): Max Clock (MC): GPU frequencies are set to maximum possible frequency domains. This is a widely used technique in the form of near-threshold computation or computational sprinting to finish the execution as fast as possible under strict performance requirements [28]. We represent our proposed policy as D-DVFS, data-driven DVFS.

Figure 7 depicts energy consumption of various applications by different policies. Both the MC and DC policies consume a much higher amount of energy as compared to our proposed data-driven frequency scaling (D-DVFS). Particularly, the MC consumes more energy than the other two policies. Since D-DVFS sensibly configures the clocks to the lowest possible energy consumption, it leads to significant energy savings for the application.

The total average GPU energy consumption can be seen in Figure 8. MC, DC, and D-DVFS have an average of 452.06 (W.S), 392.02 (W.S), and 338.01 (W.S), respectively. In other words, D-DVFS consumes 15.07% and 25.3% less than MC and DC policies, respectively. The results confirms that D-DVFS selects energy-efficient frequencies for application execution.

The application’s arrival and deadline time generated using distribution are shown in Figure 9. Accordingly, the normalized application completion time achieved using our scheduling and baseline policies is shown in Figure 10. The D-DVFS policy meets all the deadlines. It tends to execute near to the deadline requirement of applications as it scales to the frequency that has high energy efficiency and predicted execution time that meets the deadline. Although DC and MC policies execute faster, their deadline-agnostic nature tends to run the applications with high frequency and thus spending more energy. Furthermore, D-DVFS with much lower frequency executes faster in a few scenarios (for application backprop and particle_float, refer Figure 10), which represents, faster execution of applications do not have high frequency all the time. Such condition usually happens when the application has significant I/O wait or dependency stalls and setting the higher frequency in such scenarios wastes more power. Nevertheless, our approach inherently learns such behavior and adapts accordingly.

The GPU frequency scaling behavior for different applications is depicted in Figure 11. Here, MC and DC always operate statically with the highest possible and default frequencies of the GPU. However, D-DVFS selects much lower frequen-
Some works have also explored techniques like power capping and scheduling [25], [22] for energy optimization using DVFS and task mapping. The authors in [25] have used empirical, analytical models to configure the CPU-GPU frequency to execute applications within a power budget. Chau et al. [22] have studied energy-efficient job scheduling in CPU-GPU systems by regulating the DVFS. The authors proposed analytical approximation algorithms with linear programming (ILP) model and introduced a heuristic algorithm to solve this problem. Although analytical models are fast, they fail to accurately estimate the intended metrics due to their sensitivity to different parameters involved in the modeling.

The static models rely mostly on the source-code or compiler level metrics to build the models. Wang et al. [15] proposed a hybrid framework for fast and accurate GPU performance estimation through source-level analysis. They used a total of 23 parameters collected from the hardware specifications, simulation traces, and the source code. Fan et al. [16] also studied predicting the energy and performance using the static source code features from several real and synthetic open-CL kernels. Although their prediction model relies on ML techniques (Support vector regression- SVR), the training data is collected from the static source code features. They use Pareto-set of frequency configurations to find the optimal scaling values between speedup and energy further. However, models built using static features perform poorly when applied to different GPU architectures as each device has a different response to the energy and execution time. Therefore, it is beneficial to build models with actual data from the real-platforms.

Machine Learning (ML) models have been used by researchers recently in GPUs DVFS management. Wu et al. [23] proposed a neural network model to estimate the scaling curve of application with different DVFS management. While their objective is tuning different hardware parameters, we instead focus more on configuring the frequency domains to facilitate the efficient DVFS for application execution. Similarly, Guerreiro et al. [12], investigated the DVFS-aware application classification to improve GPU efficiency. They characterize the applications using the nine different application profiling features and classify the workloads based on the hierarchical clustering and neural network classifier. Our approach is different where we predict the energy and time with varying settings of the clock while this method classifies application into different domains and optimize accordingly.

Furthermore, Tang et al. [32] carried out an empirical study of GPU DVFS on energy and performance of deep learning workloads. They analyze the effect of DVFS with different core frequencies while training the deep neural networks.
on NVIDIA GT X2080 Ti. The empirical results have shown that optimal frequency settings can significantly save energy consumption. Most of these works focus on a single objective. However, in this work, we propose a data-driven frequency scaling approach for the deadline-aware scheduling algorithm.

VII. CONCLUSIONS AND FUTURE WORK

Optimal configurations of GPU frequencies can significantly reduce energy consumption. However, identifying the suitable frequencies that result in lower energy consumption with the strict application’s deadline requirement is extremely challenging. This is mainly due to the complexity induced by the application’s response to energy, execution time, and clock settings. To that end, we present a framework that selects suitable GPU frequencies for a given application using the data-driven techniques and accordingly schedule the applications while reducing energy consumption and meeting deadline. Our model achieves high accuracy with average RMSE values of 0.38 and 0.05 for energy and time, respectively, indicating that predicting the energy is quite difficult as compared to the execution time. Additionally, our proposed scheduling algorithm consumes 15.07% less energy as compared to the baselines while satisfying the deadline requirements.

In the future, we plan to extend this framework to multi-GPU platforms. We also intend to consider deep learning workloads where frequency tuning can achieve more benefits due to their long-running nature.

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